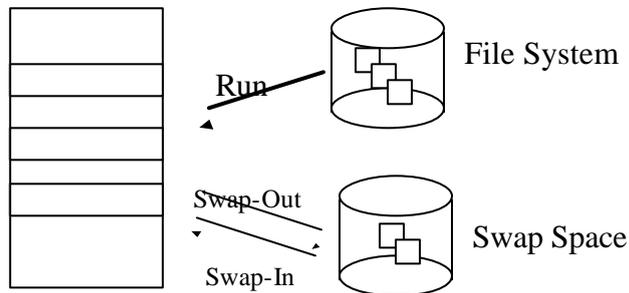


UNIX

- ✍ Introduction
- ✍ Programmer Interface
- ✍ User Interface
- ✍ Process Management
-  ✍ Memory Management
- ✍ File System
- ✍ I/O System
- ✍ Interprocess Communication

Memory Management

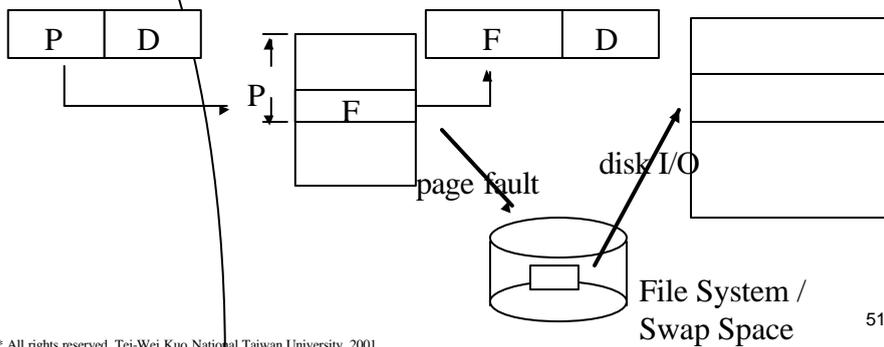
- ✍ Virtual Memory – Demand paging



Memory Management

☞ Demand Paging

- ☞ Page fault -> disk I/O -> modify the page table -> rerun the instruction!



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Memory Management

☞ 4.3BSD

☞ Reasons for page fault

☞ Initialized data and text

- ☞ Corresponding page-table entries are marked as fetch-on-demand (fod)

☞ Un-initialized data

- ☞ Corresponding page-table entries are marked as invalid -> zero-filled

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Memory Management

☞ Page Replacement

- ☞ When the number of free pages is under a threshold, some pages are paged out to release space so that the needed pages can be moved in from disks.
- ☞ Done by `pageout()` – system process `pagedaemon` is waked up when the number of free memory is less than *lotsfree*, e.g., $\frac{1}{4}$ MM size!
- ☞ Approximate Least-Recently-Used (LRU) policy

Memory Management

☞ Working Set Model

- ☞ Each process in memory should be allocated with at least those pages in the working set to prevent trashing.

☞ Swapping

- ☞ It is invoked only when paging is unable to keep up with memory needs.
- ☞ `pagedaemon` – waked up when free memory is under *minfree*, e.g., $\frac{1}{16}$ MM size.

Memory Management

✍ 3BSD

- ✍ Virtual memory – demand paging
- ✍ Page replacement – approx. LRU
- ✍ Pre-allocated swap area

✍ 4.1BSD

- ✍ Logical page – cluster
- ✍ pre-paging
- ✍ Caching recently used text pages

✍ 4.3BSD

- ✍ Text images and page tables retained in cache after exit